

SYSCALLS.1

PSIONICS FILE - SYSCALLS.1

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System calls (part 1)

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The Psion System 3 operating system offers a large number of system calls.

A system call is identified by either a "function number" or by a "functionnumber and a "subfunction number". These are abbreviated to "Fn" and "Sub" in the descriptions of system calls.

Each call can take up to 5 word parameters, called BX, CX, DX, SI, and DI. In addition, there may be up to two byte parameters (AH and AL) or an additional word parameter (AX). The values returned from the call have the same structure; in addition, there is an error flag and three result flags

(the latter are only used by a few system calls). Note that AX is equivalent to $(AH * 256 + AL)$; the description will use whichever is most convenient.

There are two OPL keywords for making system calls: CALL and OS. They access the same system calls, and differ only in the way in which the parameters are passed and the results returned.

The CALL keyword takes from 1 to 6 arguments. If an argument is omitted, the corresponding parameter is set to an undefined value. The first argument should be one of:

Fn

$Fn + 256 * Sub$

$Fn + 256 * AH$

according to the specific system call. The remaining arguments provide the word parameters:

argument 2: BX

argument 3: CX

argument 4: DX

argument 5: SI

argument 6: DI

The keyword returns a value, which is one of the results:

AX

AH * 256 + AL

(which are equivalent). The flags and the other 5 results are not available.

OS takes two or three arguments; if there is no third, it is taken to be identical to the second.

The first argument is Fn, while the second and third are each the address of a 12 byte buffer, holding the parameters and results

respectively:

Offset 0 (word): one of:

Sub * 256 + AL (parameters only)

AX (parameters and results)

AH * 256 + AL (parameters and results)

Offset 2 (word): BX

Offset 4 (word): CX

Offset 6 (word): DX

Offset 8 (word): SI

Offset 10 (word): DI The keyword returns a value which represents the error or result flags.

Each flag is held in a specific bit of the value; the remaining bits are

unspecified:

Bit 0: UL flag or error flag (depending on context)

Bit 6: EQ flag

Bit 7: SL flag

Thus, if the call can fail, it has failed if the returned value is odd.

When a parameter or result is described as a cstr, a qstr, a buffer, or any other construct requiring more than 2 bytes, the actual parameter or result is the address of the construct.

When a register holds an address (for the above reason or otherwise), it is relative to some segment register. If programming entirely in OPL, all the relevant segment registers (DS, SS, and ES) are the same. However, if an assembler is being used, it is sometimes useful to have ES different from the others (see the Psionics file PROCESS for more details). A prefix of "e" to a register

name (such as "eBX") means that the address is relative to ES. Unprefixed names are relative to DS or SS (which should be the same). A prefix of "u" means that the segment register is unknown.

The system calls are described in a standard form. The first line gives the following information:

* the Fn value

* the Sub value if any

The second line gives the following:

* The name given to the call by the SDK.* A note such as "v3.1" giving the first version of the operating system

to support the call; if omitted, all versions support it.* The word "fails" if the call can fail; if omitted, the call never fails. - If a call can fail, then the error flag shows whether or not it did so. If it failed, AL holds an error code between 128 and 255; this is normally

treated as a negative number between -1 and -128, found by evaluating the expression (AX OR \$\$\$F00). AH, and any of the 5 word results that the call would have changed, are unspecified.

- If a call cannot fail, the error flag may still be set; this should be ignored.

- Even if a call cannot fail, bad parameters can cause a system panic.* The word "async" if the call is always asynchronous, or the word "vsync" if it is sometimes asynchronous; if omitted, the call is synchronous. - A synchronous call is "completed" (has carried out all its actions) when

the system call returns. - An asynchronous call returns immediately, even though the call might not

have completed. A parameter of the call specifies a "status word", which the call initially sets to -46 (the error code meaning "not completed"). When the relevant action has been taken or event occurs, the call is completed, the status word is changed to indicate the result of the

call, and an IOSIGNAL is sent to the process. [This is exactly the same behaviour as the IOA keyword.]

- If a call is "vsync", then it may be made as either a synchronous or asynchronous call:

+ to make a synchronous call, set AL to zero (the parameter giving the address of the status word will be ignored); + to make an asynchronous call, set AL to any other value (the address

of the status word must be valid); + descriptions of such calls will omit AL from the list of parameters.

- An asynchronous call may complete before returning (for example if the relevant event has already happened). In this case the process will never see the status word with the value -46, but an IOSIGNAL is still received. - Some asynchronous calls have a corresponding "cancel" call. This causes

the original call to be completed with the status word set to some negative value other than -46.

This is then followed by a list of the parameters and results used by the call; any parameter not

listed is ignored. For the 5 main word results, the result (if the call does not fail) is the same as the corresponding parameter (including when the former is ignored) unless a different meaning is shown following an arrow (->). For the results AH, AL, and AX, and also the result flags, the result is unspecified if no such description is given.

For example, consider the following (non-existent) system call:

Fn \$12 Sub \$34

GetMagicCookie fails

AX: -> magic cookie table entry zero

BX: table entry number

CX: -> magic cookie entry

SI: addend -> CX out + SI in

Returns entries from the magic cookie table.

This system call could be used either by:

entry0% = CALL (\$3412, entry%, 0, 0, addend%) which only returns magic cookie table entry zero, or by:

LOCAL regs%(6)

regs%(1)=\$3400

regs%(2)=entryno%

regs%(5)=addend%

failed%=1 AND OS (\$12, ADDR(regs%()))

REM regs%(2,4,6) have not changed

IF NOT failed

entry0%=regs%(1)

entry%=regs%(3)

REM regs%(5) will equal entry%+addend%

ELSE

REM regs%(1,3,5) are undefined

error%=regs%(1) and \$FF00

ENDIF

which returns much more information.

System calls

Note that some system calls are deliberately not described.

Fn \$80 Sub \$00

SegFreeMemory

AX: -> free memory

Returns the amount of free system memory, in units of 16 bytes.

Fn \$80 Sub \$01

SegCreate fails

AL: segment type: ordinary processes should use 1.

AX: -> segment handle

EBX: (cstr) name of the segment

CX: segment size in units of 16 bytes
Creates an additional memory segment. Each segment must be given a name, of the form "8.3" (i.e. up to 8 characters, optionally followed by a dot and up to 3

more characters). If the name is already in use, the call will fail. Once created, the segment may be used by the process as additional memory, either via SegCopyFrom and SegCopyTo, or in assembler code (see the Psionics file

KERNEL).

Fn \$80 Sub \$02

SegDelete fails

EBX: segment handle
Deletes an additional memory segment. If any other process has opened the segment, the call will fail.

Fn \$80 Sub \$03

SegOpen fails

AX: -> segment handle

EBX: (cstr) name of the segment
Opens the additional memory segment with the given name (if no such segment

exists, the call will fail). The call will fail if the process has the segment open already.

Except where stated, all calls using additional memory segments must be given handles from SegCreate or SegOpen calls.

Fn \$80 Sub \$04

SegClose fails

BX: segment handle
Closes an additional memory segment which the process has open. If this was the

only process with the segment open, it is deleted. Open segments are closed when a process terminates.

Fn \$80 Sub \$05

SegSize

AX: -> segment size

BX: segment handle

Returns the size of an additional memory segment in units of 16 bytes.

Fn \$80 Sub \$06

SegAdjustSize fails

BX: segment handle

CX: new size in units of 6 bytes
Changes the size of an additional memory segment. The memory will be added to or removed from the end of the segment.

Fn \$80 Sub \$07

SegFind fails

AX: -> segment handle

BX: last segment handle

SI: address of 14 byte buffer used by the call

eDI: (cstr) pattern to search for
Finds the additional memory segments with names matching the search pattern (use ? and * as wild cards). The first call should be made with the lastsegment handle set to 0, and subsequent calls should set it to the handlereturned by the previous call. The call fails when there are no more segments to return. It is not necessary to keep calling SegFind until this happens.

Fn \$80 Sub \$08

SegCopyTo

BX: segment handle

CX: number of bytes to copy

DX: high half of start address in segment

SI: address in process of first byte to copy DI: low half of start address in segment

Copies data from the current process to an additional memory segment.

Fn \$80 Sub \$09

SegCopyFrom

BX: segment handle

CX: number of bytes to copy

DX: high half of start address in segment

SI: address in process to copy first byte to DI: low half of start address in segment

Copies data from an additional memory segment to the current process.

Fn \$80 Sub \$0A

SegLock

BX: segment handle
Locks an additional memory segment; a locked segment will not be deleted even if no processes have it open.

Fn \$80 Sub \$0B

SegUnLock

BX: segment handle
Unlocks an additional memory segment. The number of unlock calls should

equal

the number of lock calls; additional unlocks may cause the segment to be deleted while still open.

Fn \$80 Sub \$0C

SegRamDiskUsed

AX: -> size of the ram disc in units of 16 bytes>Returns the size of the ram disc; this should be treated carefully, as it can change without warning.

Fn \$80 Sub \$0D

SegCloseLockedOrDevice

BX: segment handleUnlocks an additional memory segment where the handle may be that for a different process. See SegUnLock for more details.

Fn \$81 Sub \$00

HeapAllocateCell fails

AX: -> allocated block

CX: number of bytes to allocateAllocates a block of memory from the heap. The returned block may be larger

than requested. Heap memory is part of the process's normal memory segment.

Fn \$81 Sub \$01

HeapReAllocateCell fails

AX: -> new address of block

BX: block to be reallocated, or zero

CX: new size of block in bytesChanges the size of a heap block; the block may have to be moved to do this.The memory will be added or deleted to the end of the block. If the call fails, the block is unaffected. If BX is zero, then a new cell is allocated. The newblock may be larger than requested.

Fn \$81 Sub \$02

HeapAdjustCellSize fails

AX: -> new address of block

BX: block to be adjusted (must not be zero)

CX: number of bytes to add or remove

DX: location within block of the changeChanges the size of a heap block; the block may have to be moved to do this.If the call fails, the block is unaffected. The memory will be added or removed

at the offset given by DX. For example, if the block was 10 bytes long, CX is3, and DX is 6, the new block will be 13 bytes, and consist of the first 6bytes of the old block, then 3 random bytes, then the last 4 bytes of the old

block. The new block may be larger than requested.

Fn \$81 Sub \$03

HeapFreeCell

BX: block to free

Frees a previously allocated heap block.

Fn \$81 Sub \$04

HeapCellSize

AX: -> cell size in bytes

BX: address of blockReturns the actual size of a heap block (which may be larger than requested

when the block was created or last resized).

Fn \$81 Sub \$05

HeapSetGranularity

BX: new granularity in units of 16 bytesWhen the application data space is not large enough to satisfy a heap request, it is grown by the granularity. The default is 2kb (128 units), and the maximum is 16k (1024 units).

Fn \$81 Sub \$06

HeapFreeMemory

AX: -> maximum possible free space on heap

BX: -> address of start of heapReturns the amount of heap space which can be allocated (by using freed blocks or by growing the data space, or both), plus the address of the base of the heap.

Fn \$82 Sub \$00

SemCreate fails

AX: -> semaphore handle

BX: initial value of semaphoreCreates a new semaphore for process interlocking. Each semaphore has a list of processes associated with it (this list is initially empty) and a value, which

must initially be zero or positive.

Fn \$82 Sub \$01

SemDelete

BX: semaphore handleDeletes a semaphore. Any processes on the semaphore's list will be restarted.

When a process terminates, any semaphore it created will be deleted.

Fn \$82 Sub \$02

SemWait

BX: semaphore handleIf the value of the semaphore is positive or zero, one is subtracted

from it.

If the value is now or initially -1, the current process is then blocked and added to the end of the list for that semaphore.

Fn \$82 Sub \$03

SemSignalOnce

BX: semaphore handle If the list for the semaphore is not empty, the first process on the list is

removed and restarted. If the list is empty (either initially or after the only process is removed), one is added to the value of the semaphore. If the restarted process has a higher priority than the current one, a reschedule will take place and the other process will start running (see the description

of priorities in the Psionics file PROCESS).

Fn \$82 Sub \$04

SemSignalMany

BX: semaphore handle

CX: number of calls to make (must be positive) This is equivalent to making several calls to SemSignalOnce.

Fn \$82 Sub \$05

SemSignalOnceNoReSched

BX: semaphore handle This is identical to SemSignalOnce except that a reschedule never takes place

because of the call (though one may take place due to the current process using up its time slot).

Fn \$83 Sub \$00

MessInit fails

BX: number of message slots + 256 * maximum data length Initialize the message system so that the current process can receive inter-process messages. The call reserves the indicated number of message slots, with each slot having room for the indicated amount of data. The slots are allocated on the heap.

Fn \$83 Sub \$01

MessReceiveAsynchronous async

BX: address of message slot pointer

DI: address of the status word When a message arrives, the message slot pointer is set to the address of the

message, the status word is set to zero, and the call completes. The message has the format:

Offset 0 to 3: used by the message system
Offset 4 (word): message type
Offset 6 (word): process sending the message
Offset 8 onward: data in message

Fn \$83 Sub \$02

MessReceiveWithWait

BX: address of message slot pointer
When a message arrives (this may have happened before the call; otherwise the call waits until a message arrives), the message slot pointer is set to the address of the message. The message format is as for MessReceiveAsynchronous.

Fn \$83 Sub \$03

MessReceiveCancel

Cancel any pending MessReceiveAsynchronous.

Fn \$83 Sub \$04

MessSend fails

BX: process to receive the message

CX: message type

SI: address of first byte of message data
Send a message to a process. The call will block until there is a free message slot in the receiving process. The receiving process determines the amount of

data sent. The reply is ignored.

If the sender has priority 128 or above, the message is placed at the front of the queue of messages waiting for the recipient (if it is not already waiting for a message). Otherwise it is placed at the back of the queue.

Fn \$83 Sub \$05

MessSendReceiveAsynchronous fails async

BX: process to receive the message

CX: message type

SI: address of first byte of message data

DI: address of the status word
Send a message to a process; when the recipient replies, the status word is set to the reply and the process is sent an IOSIGNAL. The call will block until there is a free message slot in the receiving process.

Fn \$83 Sub \$06

MessSendReceiveWithWait fails

AX: -> reply

BX: process to receive the message

CX: message type

SI: address of first byte of message data
Send a message to a process, and blocks until the

recipient replies.

Fn \$83 Sub \$07

MessFree

BX: address of received message

CX: reply

The reply is sent to the recipient of the message, and the message slot is freed and can be used for another incoming message.

Fn \$83 Sub \$08

MessSignal fails

BX: process to watch

CX: message type Requests that, when the specified process terminates, the kernel sends a message of the indicated type to the current process. The data in the message has the format:

Offset 0 (word): process id of terminating process Offset 2 (byte): reason code (see ProcKill) or panic code

Offset 3 (byte):

0 = process terminated or was killed

1 = process caused a panic

2 = a subtask of the process caused a panic in the process If the message slots do not have room for 4 bytes of data, not all this information is available.

Fn \$83 Sub \$09

MessSignalCancel fails

Fn \$83 Sub \$0A

MessSignalCancelX fails

BX: process to watch

CX: message type (MessSignalCancelX only) Cancels a call to MessSignal for the indicated process. MessSignalCancel ignores the message type and should not be used if more than one MessSignal call has been made for the process.

Fn \$84 is used to control dynamic libraries and is not described here.

Fn \$85 Sub \$00

IoOpen fails

AX: -> handle of opened device

eBX: (cstr) name of device driver

CX: mode

DX: handle of device being attached to opens a channel to a device driver. If the driver is a

base driver, then

DX is ignored (and this call is equivalent to the IOOPEN keyword). If it is a stacked driver (see the Psionics file DEVICES), then the handle to be stacked on must be specified. The meaning of the mode is determined by the device. A driver can support several channels at once, and each has its own

handle. The driver must be a logical device driver.

Fn \$85 Sub \$01 and \$02 involve physical device drivers and should only be called from within logical device drivers.

Fn \$85 Sub \$03 to \$05 should only be used by the operating system.

Fn \$85 Sub \$06

DevLoadLDD fails

Fn \$85 Sub \$07

DevLoadPDD fails

eBX: (cstr) name of file holding the driver Loads a device driver into the system. A driver cannot be opened until it is

loaded. The correct call for the driver type must be used.

Fn \$85 Sub \$08

DevDelete fails

eBX: (cstr) name of the device driver DX: \$DD01 for logical device drivers, or \$DD21 for physical device drivers Unloads a device driver from the system. Open drivers and those in the ROM cannot be unloaded.

Fn \$85 Sub \$09

DevQueryUnits fails

AX: -> number of channels supported eBX: (cstr) name of the device driver, without a trailing colon

Returns the number of simultaneous open channels supported by a driver (which must be a logical one); \$FFFF indicates no limit.

Fn \$85 Sub \$0A

DevFind fails

AX: -> find code

BX: last find code DX: \$DD01 for logical device drivers, or \$DD21 for physical device drivers

SI: address of 14 byte buffer

eDI: (cstr) pattern to search for Finds all devices with names matching the search pattern (use ? and * as wildcards). The first call should be made with the last find code set to 0, and subsequent calls should set it to the code returned by the previous call. The call fails when there are no more drivers to return. It is not necessary to keep calling DevFind until this

happens. The names do not have a trailing colon.

Fn \$85 Sub \$0B should only be used by the operating system.

Fn \$85 Sub \$0C is used for special operations on device drivers.

Fn \$86 Sub \$00

IoAsynchronous fails async

Fn \$86 Sub \$01

IoAsynchronousNoError async

Fn \$86 Sub \$02

IoWithWait fails

AL: service number

AX: -> result from driver

BX: handle of driver

CX: address of first argument

DX: address of second argument DI: address of the status word (ignored by IoWithWait)

These calls are equivalent to the IOA, IOC, and IOW keywords respectively[For those without documentation of IOC, this is the same as IOA, except that errors are handled by setting the status word and calling IOSIGNAL, so that

IOC always succeeds, unlike IOA.]

Fn \$86 Sub \$03 and \$04 should only be used by device drivers.

Fn \$86 Sub \$05

IoWaitForSignal

This call is equivalent to the IOWAIT keyword.

Fn \$86 Sub \$06

IoWaitForStatus

DI: address of the status word

This call is equivalent to the IOWAIT keyword.

Fn \$86 Sub \$07

IoYield

DI: address of the status word

This call is equivalent to the IOYIELD keyword.

Fn \$86 Sub \$08

IoSignal

Sends an IOSIGNAL to the current process. This call is equivalent to the IOSIGNAL keyword.

Fn \$86 Sub \$09

IoSignalByPid fails

Fn \$86 Sub \$0A

IoSignalByPidNoReSched fails

BX: process ID Sends an IOSIGNAL to a process. With the latter call a reschedule never takes place because of the call (though one may take place due to the current process using up its time slot).

Fn \$86 Sub \$0B to \$0F should only be used by device drivers.

Fn \$86 Sub \$10

IoClose fails

AX: -> result from driver

BX: handle

This call is equivalent to the IOCLOSE keyword.

Fn \$86 Sub \$11

IoRead fails

Fn \$86 Sub \$12

IoWrite fails

AX: -> amount actually read or result from driver

BX: handle

CX: address of buffer

DX: number of bytes read or written These calls are equivalent to the IOREAD and IOWRITE keywords.

Fn \$86 Sub \$13

IoSeek fails

AX: -> result from driver

BX: handle

CX: mode

DX: address of long holding seek position

This call is equivalent to the IOSEEK keyword.

Fn \$86 Sub \$14 should only be used by the window manager.

Fn \$86 Sub \$15 to \$17 should only be used from assembler.

Fn \$86 Sub \$18

IoShiftStates

AL: modifiers

This call makes available the state of the various modifier keys:

Bit 1: shift

Bit 2: control

Bit 3: psion

Fn \$86 Sub \$19

IoWaitForSignalNoHandlerThis call should be used instead of IOSEEK by tasks (subsidiary processes of a process).

Fn \$86 Sub \$1A

IoSignalKillAsynchronous fails async

BX: process to watch

DI: address of the status wordThis call completes, and the status word is set to 0, when the specified process terminates.

Fn \$86 Sub \$1B

IoSignalKillCancel fails

BX: process to watch

Cancel any pending IoSignalKillAsynchronous.

Fn \$86 Sub \$1C and \$1D should only be used by the window manager.

Fn \$86 Sub \$1E

IoPlaySoundW v3

AL: -> failure code

BX: (cstr) sound file name

CX: duration to play (in 1/32 second)

DX: volume: 0 (loudest) to 5 (softest)Plays a sound file. A duration of 0 means the file header specifies the duration; otherwise, the sound is clipped or padded with space as needed.The name may be either an ordinary filename, or a string of the form "*ABCD", meaning the first found of the files:

ROM::ABCD.WVE

LOC::M:\WVE\ABCD.WVE

LOC::A:\WVE\ABCD.WVE

LOC::B:\WVE\ABCD.WVENote that this call does not fail in the usual way, but just returns a failure code.

Fn \$86 Sub \$1F

IoPlaySoundA v3 async

BX: (cstr) sound file name

CX: duration to play (in 1/32 second)

DX: volume: 0 (loudest) to 5 (softest)

DI: address of the status wordPlays a sound file asynchronously; the call completes when the sound has

finished. The other arguments are as for IoPlaySoundW.

Fn \$86 Sub \$20

IoPlaySoundCancel v3

Cancel any pending IoPlaySoundA.

Fn \$86 Sub \$21

IoRecordSoundW v3

AL: -> failure code

BX: (cstr) sound file name CX: number of samples to record, in units of 2048 samples

Records a sound file. Note that 2048 samples are slightly more than a quarter of a second. The file will be created before recording, and must not be on a flash device. Note that this call does not fail in the usual way, but just returns a failure code.

Fn \$86 Sub \$22

IoRecordSoundA v3 async

BX: (cstr) sound file name

CX: number of samples to record, in units of 2048 samples DI: address of the status word

Records a sound file asynchronously; the call completes when the recording has finished. The other arguments are as for IoRecordSoundW.

Fn \$86 Sub \$23

IoRecordSoundCancel v3

Cancel any pending IoRecordSoundA.

Fn \$86 Sub \$24

IoPlaySoundA0 v3.9 async

BX: (cstr) sound file name

CX: duration to play (in 1/32 second)

DX: volume: 0 (loudest) to 5 (softest)

DI: address of the status word

SI: duration to skip (in 1/32 second) Plays part of a sound file asynchronously, skipping some initial portion of the sound; the call completes when the sound has finished. The other arguments are as for IoPlaySoundA.

Fn \$87 Sub \$00 is carried out automatically for OPL programs.

Fn \$87 Sub \$01

FileExecute fails vsync

eBX: (cstr) program file name

CX: (qstr) program information

DX: address of the status word

DI: address of word to be filled with process ID of new

process

Starts a new process and places it in "suspended" state. The program filename will have the extension ".IMG" added if one is not specified. The program information is a qstr whose interpretation depends on the program; it is available from location 36 of the new process's data segment (see Psionics file PROCESS). The process name will be set to the name part of the file (excluding any extension). If another process of this name is already running, both must execute the same code, or the call will fail.

@@ More in S3/S3a manual chapters 1 and 2. @@ Program information often consists of the following catenated

together (note that all these elements except the first are cstrs): (1) the character 'C' (create new file) or 'O' (open existing file); (2) (cstr) the application name, with the first character uppercased and the

rest lowercased;

(3) (cstr) one of:

(a) the normal extension for filenames, including the leading dot; (b) the normal extension with dot, followed by a space and aliasing

information;

(c) an empty string

(4) (cstr) the data file path passed as the initial argument; (5) (cstr) the application path (may be omitted if not required).

A binary application may be used to operate on a standard file type. Examples include:

Program	Application	Extension	[item 2]
"ROM::DATA.APP"	"Data"	".DBF"	"ROM::WORD.APP"
"ROM::WORD.APP"	"Program"	".OPL"	"Word"
"LOC::C:\APP\COMMS.APP"	"Comms"	".SCO"	".WRD"

In each of these cases there is no aliasing information or application path. An example program information qstr is "OData~.DBF~LOC::M:\DBF\MYDATA.DBF~" (tilde indicates a zero byte).

@This section needs rewriting@The word processor may also be used as a program editor, by making item 2

its name ("Program") and item 3 an extension (usually ".OPL") followed by the necessary aliasing information: this is empty for normal Word, or: * "O" for OPL editor, "S" for SCR editor, "\$" for plain text editor (3a only)

"/" for custom template (3a only); any letter uses translator SYS\$PRG<letter>.IMG to

translate, then

* then "R" indicates that the translator can execute resulting code, anything else means it can't, then

* extension and directory ("OPO" or "SCO" for example) of translated files,* then on 3a only optional "*" to indicate compatibility mode exists. An example program information qstr is "Program~.OPL OROPO~LOC::M:\OPL\MYPROG.OPL~"

template mode for Word means there must be a \WDR\

template name is that if the application name. For example,

alias file of:

Letter.LET

\LET\

1083

Word

/

uses template LETTER. Note mode 80 in application type in alias file.

Finally, a translator may be used to execute a translated file: in the case of the translator "ROM::SYS\$PRG0", it can execute either OPO or OPA files. The application name is given in the APP ... ENDA clause for OPA files, and is "RunOpl" for OPO files; there is no extension or aliasing information.

For OPA files, the data file is passed to the application, and the application path shows which application is actually run. For OPO files, the data file should name the OPO file itself, and there is no application path. Examples are "RunOpl~M:\OPO\MYPROG.OPO~" and "MyApp~M:\DAT\MYDATA~M:\APP\MYAPP.OPA~".

Fn \$87 Sub \$02

FileParse fails vsync

BX: (cstr) file name to be parsed

CX: (cstr) file specification to be used

DX: address of the status word

SI: address of a 6 byte buffer to be filled in DI: address of 128 byte buffer to be filled with the parsed filename This call is equivalent to the PARSE\$ keyword. The buffer is filled as follows:

Offset 0 (byte): length of node name

Offset 1 (byte): length of device name

Offset 2 (byte): length of path name

Offset 3 (byte): length of base name

Offset 4 (byte): length of extension

Offset 5 (byte): flags Bit 0: the base name or extension contains a wildcard ("*" or

"?")

Bit 1: the base name contains a wildcard

Bit 2: the extension contains a wildcard

Fn \$87 Sub \$03

FilPathGet fails vsync

BX: 128 byte buffer

DX: address of the status wordThe buffer is filled with the current filing system default path (a cstr).

Fn \$87 Sub \$04

FilPathSet fails vsync

BX: (cstr) new path

DX: address of the status word

Sets the default path (equivalent to the SETPATH keyword).

Fn \$87 Sub \$05

FilPathTest fails vsync

BX: (cstr) pathname to be tested

DX: address of the status wordEquivalent to the EXIST keyword; the call succeeds if the file with that pathname exists.

Fn \$87 Sub \$06

FilDelete fails vsync

BX: (cstr) pathname to be deleted

DX: address of the status wordEquivalent to the DELETE keyword; non-empty directories cannot be deleted.

Fn \$87 Sub \$07

FilRename fails vsync

BX: (cstr) old pathname

CX: (cstr) new pathname

DX: address of the status word

Equivalent to the RENAME keyword.

Fn \$87 Sub \$08

FilStatusGet fails async

BX: (cstr) filename

CX: 16 byte buffer

DX: address of the status word

The buffer is filled in with information about the file: Offset 0 (word): buffer format version: always 2

Offset 2 (word): attributes
Offset 4 (long): file size in bytes
Offset 8 (long): time last modified

The attribute bits have the following meanings (question marks indicate that the bit name was specified but the meaning is unknown; equals signs mean that the bit can be set by

FileStatusSet):

= Bit 0: file is not ReadOnly
= Bit 1: file is Hidden
= Bit 2: file is a System file
Bit 3: file is a Volume name
Bit 4: file is a Directory
= Bit 5: file is Modified
? Bit 8: file is a readable file
? Bit 9: file is an executable file
? Bit 10: file is a byte stream file

Bit 11: file is a text file
The LOC:: node (see Psionics file FILEIO) does not distinguish text and binary files, and always leaves bit 11 clear. Other nodes may distinguish the types of files and set the bit where appropriate.

Fn \$87 Sub \$09

FileStatusSet fails vsync

BX: (cstr) filename

CX: mask of attributes to be changed (only specify settable bits) DX: address of the status word

DI: new values for attributes to be altered
The attributes indicated in the mask of the specified file are altered to the new values; all other attributes are left unchanged.

Fn \$87 Sub \$0A AL sync

FileStatusDevice fails vsync

BX: (cstr) device name (such as "A:" or "LOC::A:")

CX: 64 byte buffer

DX: address of the status word

The buffer is filled in with information about the device: Offset 0 (word): buffer format version

Offset 2 (byte): device type

0 = unknown

1 = floppy

2 = hard disc

3 = flash

4 = ram

5 = rom

6 = write-protected

Offset 3 (byte): device properties:

Bit 3: formattable device

Bit 4: dual density device

Bit 5: internal device

Bit 6: dynamically sizeable device Bit 7: compressible (worth compressing database

files)

Offset 4 (word): non-zero if the device contains removable media Offset 6 (long): total space on the device in bytes

Offset 10 (long): free space on the device in bytes

Offset 14 (cstr): volume name

Offset 46 (word): device battery status

0 = low battery

1 = battery OK

-4 = status not available for this device type This field is not valid if the format version (offset 0) is less than 3.

Fn \$87 Sub \$0B

FilStatusSystem fails vsync

BX: (cstr) node name (such as "LOC:." or "REM:.")

CX: 32 byte buffer

DX: address of the status word

The buffer is filled in with information about the node: Offset 0 (word): buffer format version (currently 2)

Offset 2 (word): 0 = flat, 1 = hierarchical Offset 4 (word): non-zero if devices are formattable, and zero otherwise

Fn \$87 Sub \$0C

FilMakeDirectory fails vsync

BX: (cstr) pathname

DX: address of the status word

This is equivalent to the MKDIR keyword.

Fn \$87 Sub \$0D AL sync

FilOpenUnique fails vsync

eBX: (cstr) pathname

CX: open mode (format and access parts only)

DX: address of the status word
IOOPEN keyword with
a mode of 4. The pathname is used to determine the directory to create the file in, and will be modified to give the actual name (thus the pathname should have room for 128 characters).

Fn \$87 Sub \$0E

FileSystemAttach fails vsync

BX: (cstr) name of a file system PDD

DX: address of the status word
The specified physical device driver will be attached to the filing system,
possible adding new nodes.

Fn \$87 Sub \$0F

FileSystemDetach fails vsync

BX: (cstr) name of a filing system

DX: address of the status word
Detaches a filing system. Built-in filing systems cannot be detached.

Fn \$87 Sub \$10

FilPathGetById fails vsync

BX: process to examine

CX: 128 byte buffer

DX: address of the status word
The buffer is set to the current path of the specified process (a cstr).

Fn \$87 Sub \$11

FilChangeDirectory fails vsync

BX: (cstr) pathname

CX: 128 byte buffer

DX: address of the status word

SI: (cstr) name of subdirectory DI: 0 = get root directory, 1 = get parent directory, 2 = get subdirectory

The path name is modified in the requested way. This call should be used instead of manipulating the directory part of a name directly, as it works on all nodes, irrespective of the format that the node uses for path names.

Fn \$87 Sub \$12 has no effect on OPL processes.

Fn \$87 Sub \$13

FilSetFileDate fails vsync

BX: (cstr) pathname

CX: low word of new time

DX: address of the status word

DI: high word of new time

Sets the modification time of the specified file.

Fn \$87 Sub \$14

FillLocChanged fails vsync

AX: -> channel change flags

BX: channel to check

DX: address of the status word Specifies whether any directory in the LOC:: node has changed (i.e. a file or directory has been created, destroyed, or renamed, but not changes to file contents) since the last use of this system call with this channel. BX should have exactly one bit set, corresponding to the channel to use; bits 8 to 15 are reserved by Psion. The corresponding bit in AX will be set if the node has changed, and will be clear otherwise; the other 15 bits are unspecified. This call can be used to determine whether to update a directory display.

Fn \$87 Sub \$15

FillLocDevice v3 fails vsync

BX: local device indicator (%A to %H, or %I or %M) CX: address of word set to the media type

DX: address of the status word Provides the media type of a device on the LOC:: node (I and M both refer to

the internal ramdisc). The device type consists of:

Bits 0 to 3:

0 = unknown

1 = floppy

2 = hard disc

3 = flash

4 = ram

5 = rom

6 = write-protected

Bits 7 to 8:

0 = device battery measurement not supported 1 = device battery measurement not supported

2 = device battery voltage low

3 = device battery voltage good

Fn \$87 Sub \$16

FillLocReadPDD v3 fails vsync

BX: local device indicator (%A to %H, or %I or %M)

CX: (long) location on the device

DX: address of the status word

SI: number of bytes to read

DI: address in process to copy first byte toCopies data from a device on the LOC:: node to the current process. This call is very efficient, and accesses the raw device.

Further system calls are described in Psionics file SYSCALLS.2.

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